

# Generalized $\beta$ -skeletons

1

**GABRIELA MAJEWSKA  
INSTITUTE OF INFORMATICS  
UNIVERSITY OF WARSAW  
POLAND**

**MIROSLAW KOWALUK  
INSTITUTE OF INFORMATICS  
UNIVERSITY OF WARSAW  
POLAND**

# $\beta$ -skeletons

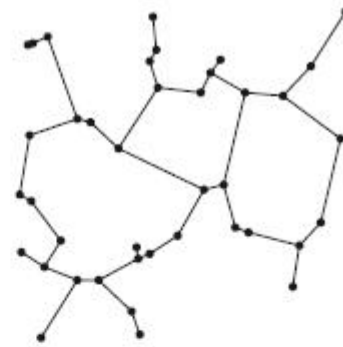
2

The  **$\beta$ -skeletons  $\{G(V)\}_\beta$**  for a point set  $V$  is a hierarchy of graphs on  $V$  based upon a natural notion of „neighborlines” parameterized by real number  $\beta \geq 0$ . They are both important and popular because of many practical applications which span a spectrum of areas from geographic information systems and wireless ad hoc networks to shape recognition and machine learning.

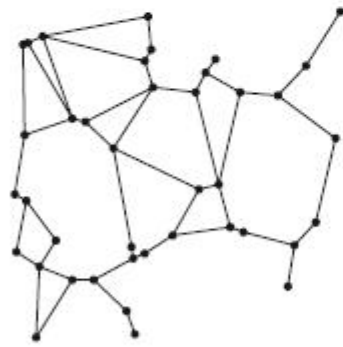
Two types of  $\beta$ -skeletons are especially well-known, Gabriel Graph (GG) for  $\beta = 1$  and the Relative Neighborhood Graph (RNG) for  $\beta = 2$ .



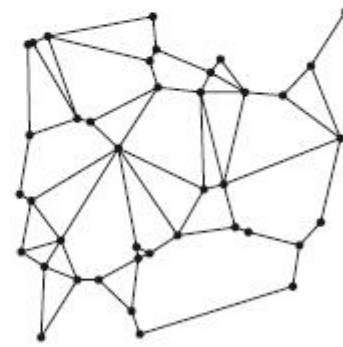
Point Set



$\beta = 2.0$ , 40 points and 42 edges



$\beta = 1.50$ , 40 points and 52 edges



$\beta = 1.0$ , 40 points and 66 edges

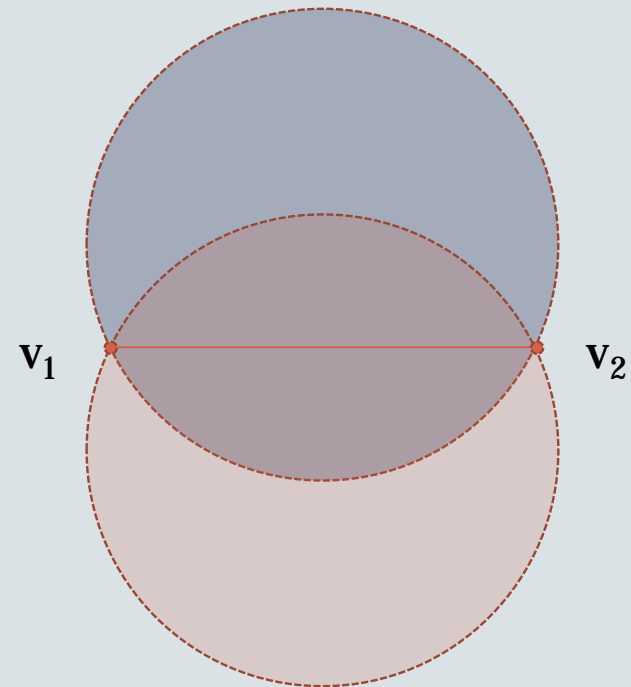
Figure 2:  $\beta$ -skeletons for various values of  $\beta$

# Lune-based $\beta$ -skeletons

4

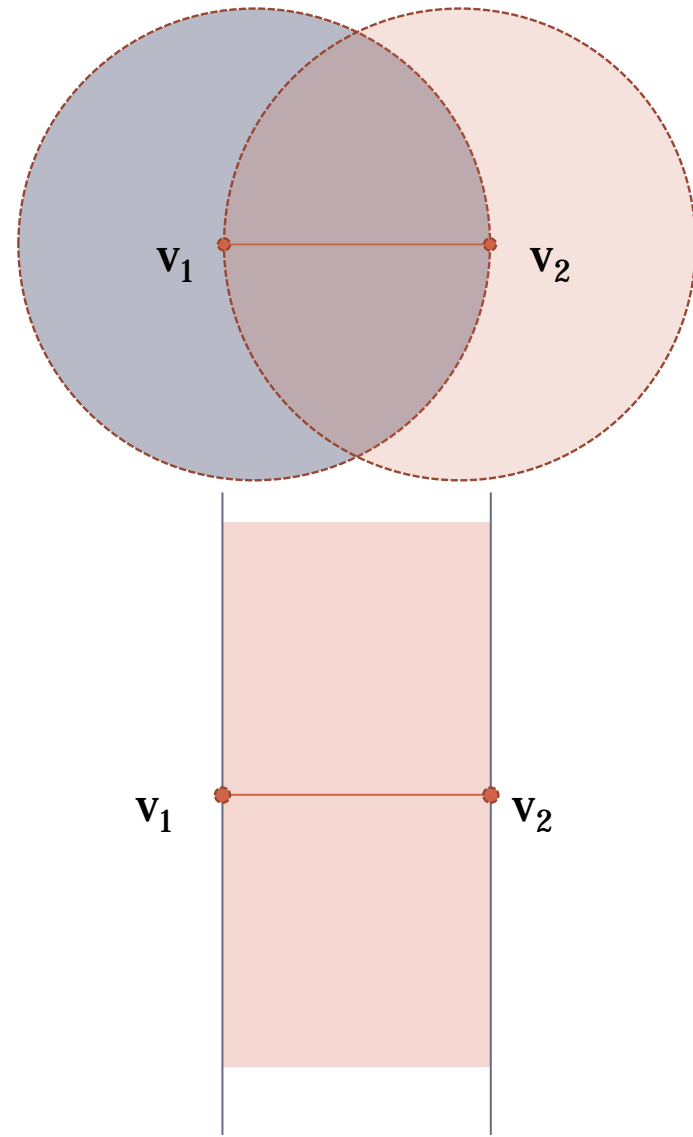
For a given set of points  $V = \{v_1, v_2, \dots, v_n\}$  in  $\mathbb{R}^2$  and parameters  $\beta \geq 0$  and  $p$  we define graph  $G_\beta(V)$  - called a *lune-based  $\beta$ -skeleton* - as follows: two points  $v_1, v_2$  are connected with an edge if and only if no point from  $V \setminus \{v_1, v_2\}$  belongs to the set  $N_p(v_1, v_2, \beta)$  where:

1. for  $\beta=0$  the set  $N_p(v_1, v_2, \beta)$  is segment  $v_1v_2$ ;
2. if  $0 < \beta < 1$  then  $N_p(v_1, v_2, \beta)$  is an intersection of two discs in  $l_p$ , each with radius  $|v_1v_2|/2\beta$ , whose boundaries contain both  $v_1$  and  $v_2$ ;



3. for  $1 \leq \beta < \infty$  set  $N_p(v_1, v_2, \beta)$  is an intersection of two discs in  $l_p$  metric, each with radius  $\beta|v_1v_2|/2$ , whose centers are in points  $(\beta/2)v_1 + (1-\beta/2)v_2$  and  $(1-\beta/2)v_1 + (\beta/2)v_2$ ;

4. for  $\beta = \infty$ ,  $N_p(v_1, v_2, \beta)$  is the unbound strip between lines perpendicular to the segment  $v_1v_2$  that contain  $v_1$  and  $v_2$  respectively.



# Motivation

6

## Fact 1

Let us assume that that points in  $V$  are in general position.

For  $1 \leq \beta \leq \beta' \leq 2$  following inclusions are true:

$$\text{MST}(V) \subseteq \text{RNG}(V) \subseteq G_{\beta'}(V) \subseteq G_{\beta}(V) \subseteq \text{GG}(V) \subseteq \text{DT}(V).$$

Notice, that in some metric spaces (like for example in weighted graphs) it is hard to interpret the linear combination of analyzed two points from the definition of the lune-based  $\beta$ -skeleton, so we want to instead use the fact that we know the distances of centers of discs defining the lune from the ends of the edge we are checking.

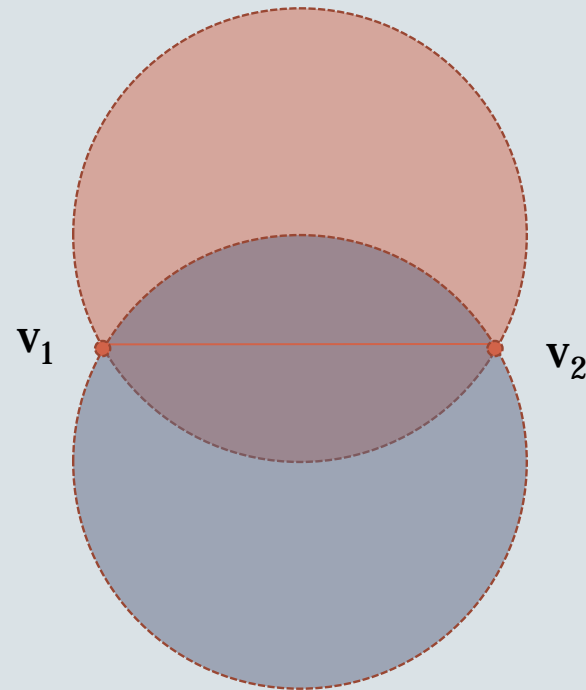
We are interested in creating a global definition, based only on a distance criterion, that will allow us to define  $\beta$ -skeletons for a bigger class of events and that will satisfy the above inclusions.

# Circle-based $\beta$ -skeletons

7

By changing the definition of lune-based  $\beta$ -skeletons for  $\beta \geq 1$  we get the family of circle-based  $\beta$ -skeletons.

1. for  $1 \leq \beta < \infty$  set  $N_p(v_1, v_2, \beta)$  is an union of two discs in  $l_p$  each with radius  $\beta |v_1 v_2|/2$ , whose boundaries contain both  $v_1$  and  $v_2$ ;
2. for  $\beta = \infty$ ,  $N_p(v_1, v_2, \beta)$  is an union of segment  $v_1 v_2$  and two open hyperplanes defined by the line passing through  $v_1 v_2$ .



# Defining $\beta$ -skeletons using a distance criterion

8

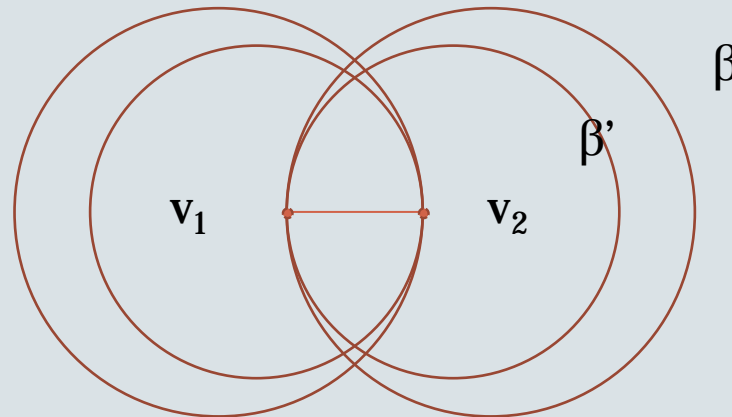
For a given set  $V$  and given parameters  $\beta \geq 0$  and  $1 < p < \infty$  we define a graph  $G_\beta^*(V)$ , where an edge exists between two points  $v_1$  and  $v_2$  iff no point from  $V \setminus \{v_1, v_2\}$  belongs to set  $N_p^*(v_1, v_2, \beta)$  where:

1. for  $0 \leq \beta < 1$  we have  $N_p^*(v_1, v_2, \beta) = N_p(v_1, v_2, \beta)$  defined as before;
2. for  $1 \leq \beta < \infty$  the set  $N_p^*(v_1, v_2, \beta)$  is an intersection of two  $l_p$  discs, each with radius  $\beta |v_1 v_2|/2$ , whose centers are in such points  $c_1$  and  $c_2$  that the distance between  $v_1$  and  $c_1$  ( $v_2$  and  $c_2$  respectively) is  $\beta |v_1 v_2|/2$ , the distance between  $v_1$  and  $c_2$  ( $v_2$  and  $c_1$  respectively) is  $|1 - \beta/2| |v_1 v_2|$  and the distance between  $c_1$  and  $c_2$  is  $(\beta - 1) |v_1 v_2|$ ;



3. If  $\beta = \infty$  we define  $C_\infty$  as a set of all discs  $c$  such that there exists a sequence of discs  $\{c(\beta) \mid c(\beta) \text{ is a disc that we used to define set } N_d(v_1, v_2, \beta)\}$  convergent to  $c$  with  $\beta \rightarrow \infty$ ;  
 then  $N_d(v_1, v_2, \infty)$  is intersection of any two different discs from  $C_\infty$ .

We rely here on the fact that for  $\beta \geq \beta' \geq 1$  lune defined by discs  $D_p(c_1(\beta'), \beta' |v_1 v_2|/2)$  and  $D_p(c_2(\beta'), \beta' |v_1 v_2|/2)$  is contained in a similar lune defined for  $\beta$ .



### **Lemma 1**

Let  $V$  be a set of points in  $\mathbb{R}^2$  with  $l_p$  metric, where  $1 < p < \infty$ . Then, centers of the discs determining lunes are uniquely defined and  $G_\beta(V) = G_\beta^*(V)$ .

Note that in  $l_1$  and  $l_\infty$  Lemma 1 is not true, since it is possible to find in those metrics internally tangent circles with different centers that intersect at many points and so the centers of the discs determining lunes are not uniquely defined.

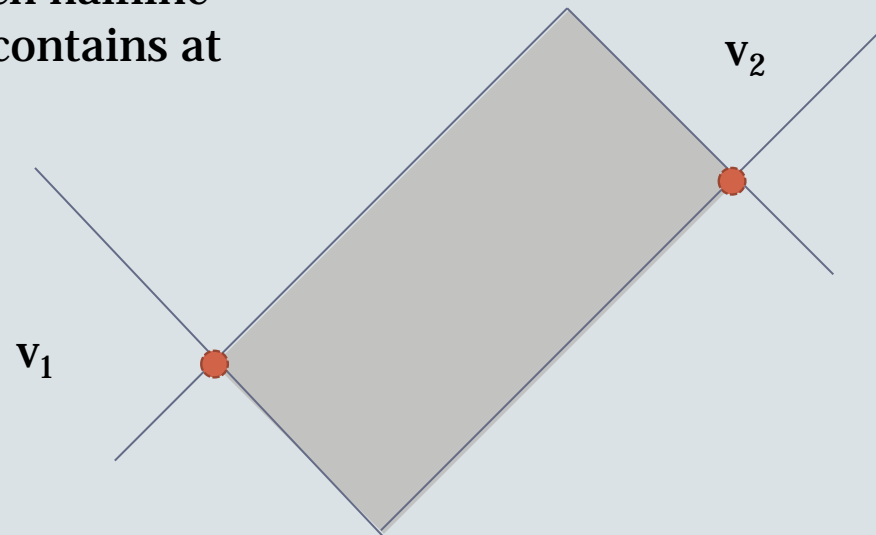
We are interested in changing the definition so that we also include the cases where the centers of the discs defining the lunes for a given edge are not uniquely defined.

# Problem 1: Not uniquely defined centers of the discs determining lunes

11

In  $l_p$  metric, where  $p \in \{1, \infty\}$ , for a given set  $V$  and given parameter  $\beta \geq 0$  we define a graph  $G_\beta^\wedge(V)$ , where an edge exists between two points  $v_1$  and  $v_2$  iff no point from  $V \setminus \{v_1, v_2\}$  belongs to one of the sets  $N_p^\wedge(v_1, v_2, \beta)$  where:

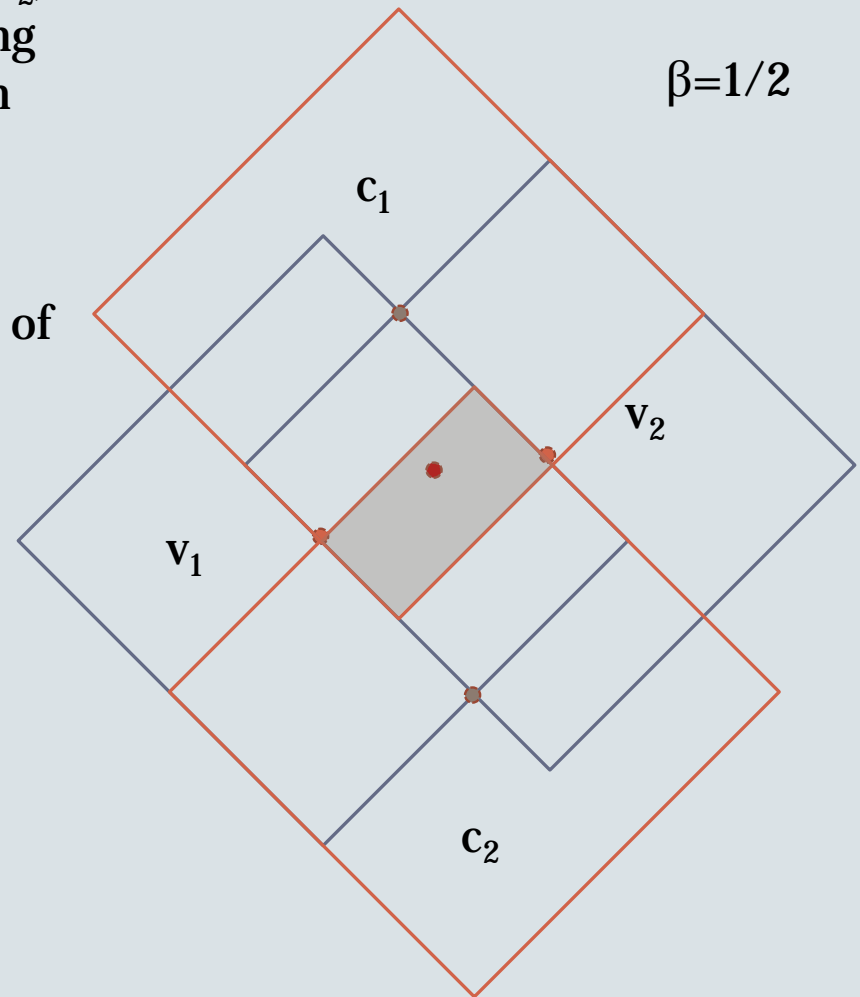
1. for  $\beta=0$  the set  $N_p^\wedge(v_1, v_2, \beta)$  is an intersection of two quarter-planes (that borders are parallel to parts of circles in  $l_p$ ), such that each halfline on the border of the quarter-planes contains at least one of points  $v_1$  and  $v_2$ ;



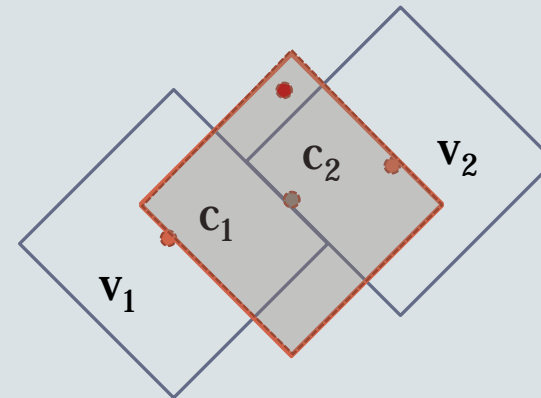
2. if for  $0 < \beta < 1$  let points  $c_1$  and  $c_2$  be such points (if they exist) that:  
 $d_p(c_1, v_1) = d_p(c_2, v_1) = d_p(c_1, v_2) = d_p(c_2, v_2) = d_p(v_1, v_2) / 2\beta$  and all paths connecting those points with length shorter than  $d_p(v_1, v_2) / \beta$  intersect shortest paths between  $v_1$  and  $v_2$ :

then  $N_p^\wedge(v_1, v_2, \beta)$  is an intersection of discs  $D_p(c_1, d_p(v_1, v_2) / 2\beta)$  and  $D_p(c_2, d_p(v_1, v_2) / 2\beta)$ ;

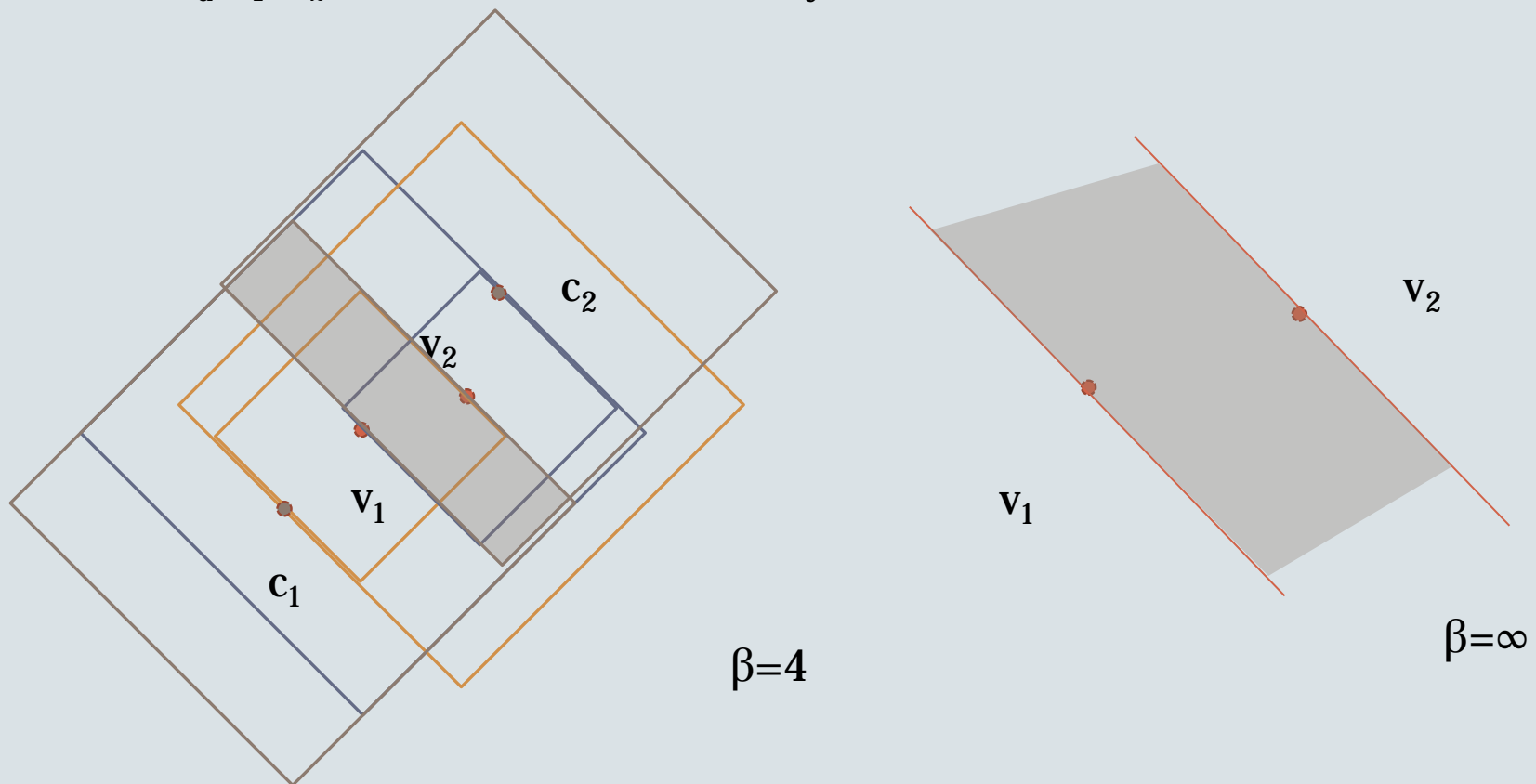
if points  $c_1$  and  $c_2$  don't exist then  $N_p^\wedge(v_1, v_2, \beta)$  is an empty set;



3. for  $1 \leq \beta < \infty$  we define set  $C_1$  (respectively  $C_2$ ) of disc centers  $c_1$  such that the distance between  $v_1$  and  $c_1$  ( $v_2$  and  $c_2$  respectively) is  $\beta |v_1 v_2| / 2$  and the distance between  $v_1$  and  $c_2$  ( $v_2$  and  $c_1$  respectively) is  $|1 - \beta/2| |v_1 v_2|$ ; set  $N_p^\wedge(v_1, v_2, \beta)$  is an intersection of any two discs, each with radius  $\beta |v_1 v_2| / 2$  centered in such points  $c_1 \in C_1$  and  $c_2 \in C_2$  (respectively) that the distance between  $c_1$  and  $c_2$  is  $(\beta - 1) |v_1 v_2|$ ;



4. for  $\beta = \infty$  we define  $C_\infty$  as a set of all discs  $c$  such that there exists a sequence of discs  $\{c(\beta) \mid c(\beta) \text{ is a disc that we used to define set } N_d(v_1, v_2, \beta)\}$  convergent to  $c$  with  $\beta \rightarrow \infty$ ;  
 then  $N_d(v_1, v_2, \infty)$  is intersection of any two different discs from  $C_\infty$ .



## Problem 2: $\beta$ -skeleton for a set of objects different than points

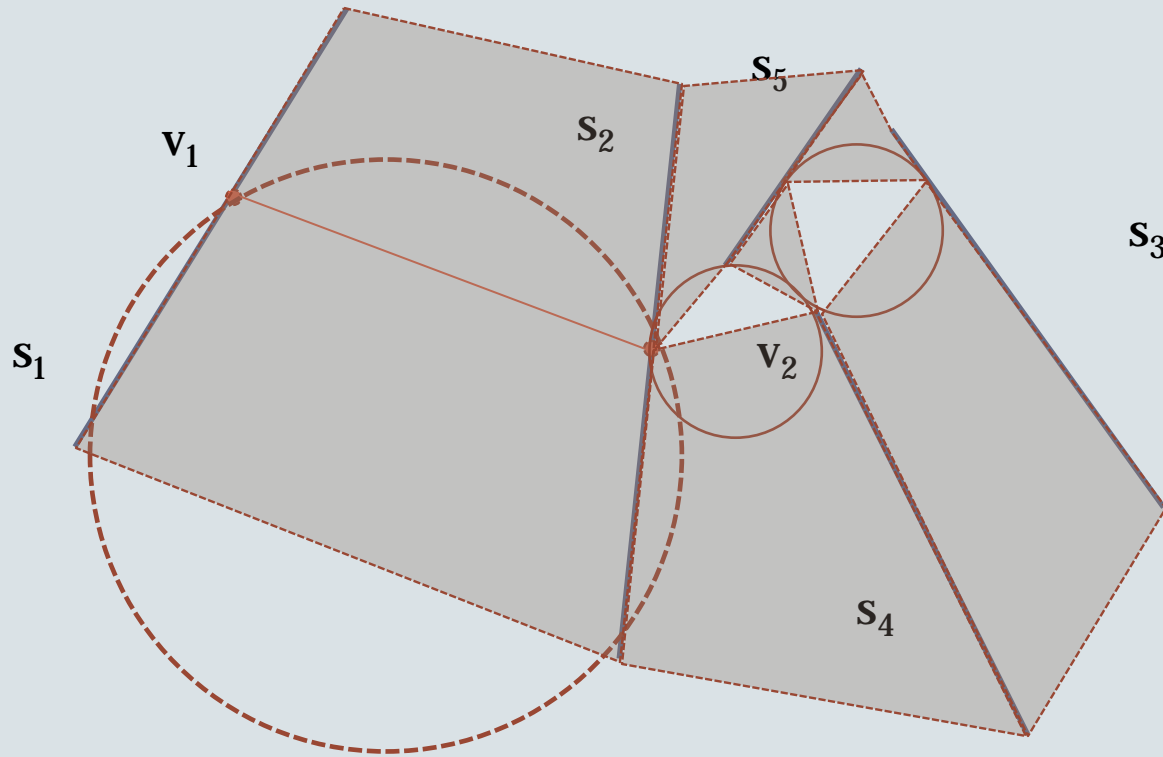
15

Let  $S$  be a set of segments in  $(\mathbb{R}^2, l_p)$ . From *Flip Algorithm for Segment Triangulations* by Brevilliers, Chevallier and Schmitt we have:

A **segment triangulation**  $T$  of  $S$  is a partition of the convex hull  $\text{conv}(S)$  of  $S$  in disjoint sites, edges, and faces such that:

1. Every face of  $T$  is an open triangle whose vertices are in three distinct sites of  $S$  and whose open edges do not intersect  $S$ ,
2. No face can be added without intersecting another one,
3. The edges of  $T$  are the (possibly two-dimensional) connected components of  $\text{conv}(S) \setminus (F \cup S)$ , where  $F$  is the union of faces of  $T$ .

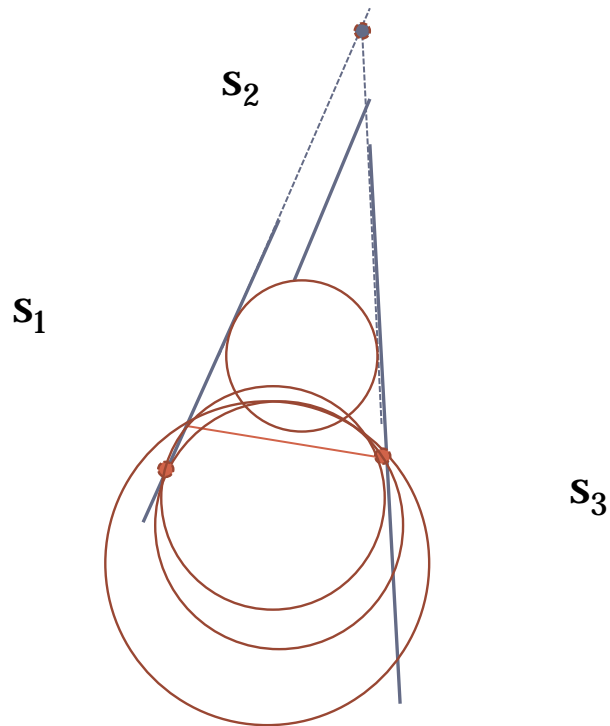
A segment triangulation of  $S$  is Delaunay  $DT(S)$  if the circumcircle of each face does not contain any point of  $S$  in its interior.



Now, for a given parameter  $\beta \in [1, 2]$  we define lune-based  $\beta$ -skeleton  $M_\beta(S)$  as follows:

for segments  $s_1$  and  $s_2$  and edge exists between vertices  $v_1 \in s_1$  and  $v_2 \in s_2$  iff no point from segments  $S \setminus \{s_1, s_2\}$  belongs to set  $N_2(v_1, v_2, \beta)$  (defined like in the first or second definition).





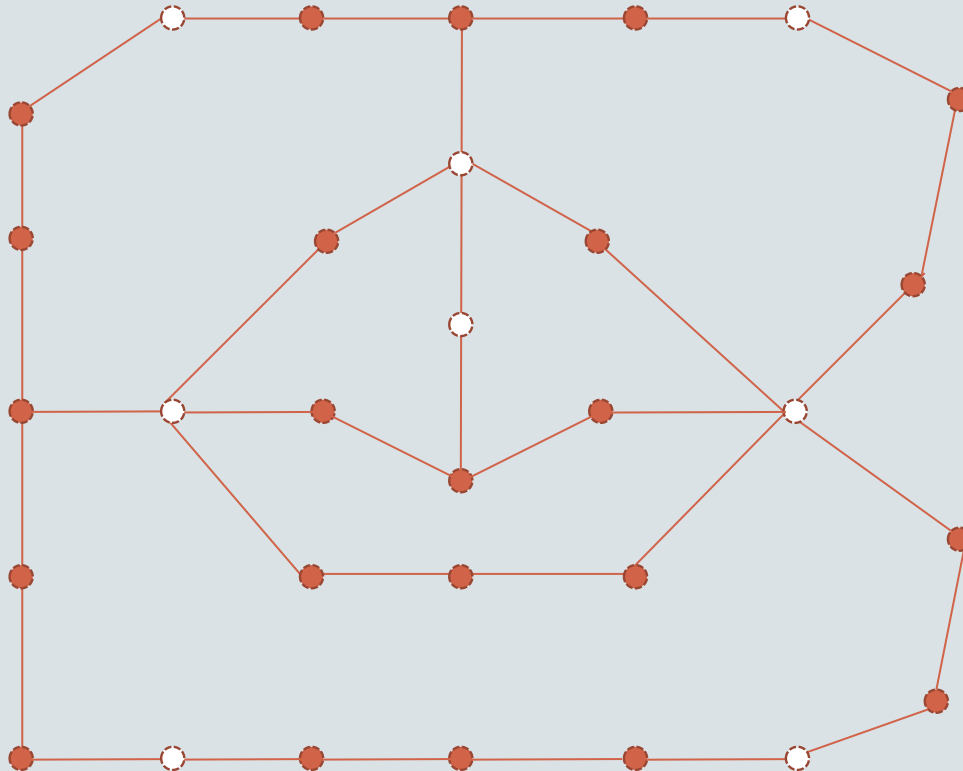
## Observation 2

For  $1 \leq \beta \leq \beta' \leq 2$  we have:

$MST(S) \subseteq RNG(S) \subseteq M_{\beta'}(S) \subseteq M_{\beta}(S) \subseteq GG(S) \subseteq DT(S)$   
 where  $RNG(S) = M_2(S)$ ,  $GG(S) = M_1(S)$  and  $MST(S)$  is a minimum spanning tree for the set of segments  $S$ .

# $\beta$ -skeletons in weighted graphs 1

18



Graph  $G=(V,U,E)$   
where all edges have a  
positive, finite weight

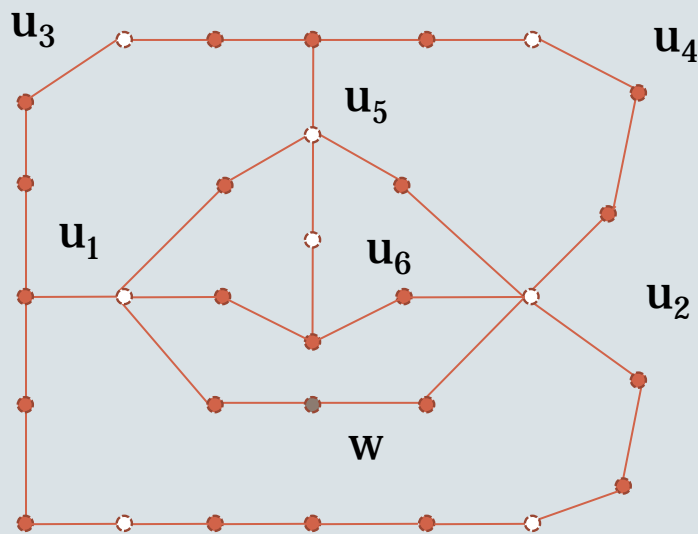
We can define a metric  
 $d_G$  here:  
 $d_G(v,w) = |\text{shortest path}$   
between  $v$  and  $w|$

From *Proximity graphs in large weighted graphs* by Abrego et al. we get:

The **Voronoi region** for point  $u_i \in U$  consists of those points  $p \in G$  that satisfy  $d_G(p, u_i) \leq d_G(p, u_k)$  for all  $u_k \in U \setminus \{u_i\}$ . **Voronoi Diagram of G**, denoted by  $VD(G)$  is the partition of  $G$  into Voronoi regions for points  $u_i \in U$ .

Let  $DT(G)$  be such a graph where points  $u_i$  and  $u_j$  are connected by an edge if and only if there exists a disc  $D_G(v, r)$  enclosing  $u_i$  and  $u_j$  and containing no other points from  $U$ , where  $v \in V$  and  $r = \min\{r \mid D_G(v, r) \text{ contains } u_i \text{ and } u_j\}$ .

We call this graph **Delaunay Triangulation** of  $V$



# $\beta$ -skeletons in weighted graphs 2

20

In graph  $G=(V,U,E)$  with metric  $d_G$  for a given parameter  $\beta \geq 0$  we define a graph  $G_\beta(G)=(U,F)$ , where an edge exists between two points  $u_1$  and  $u_2$  iff there exists a set  $N_G(u_1, u_2, \beta)$  such that no point from  $U \setminus \{u_1, u_2\}$  belongs to it where:

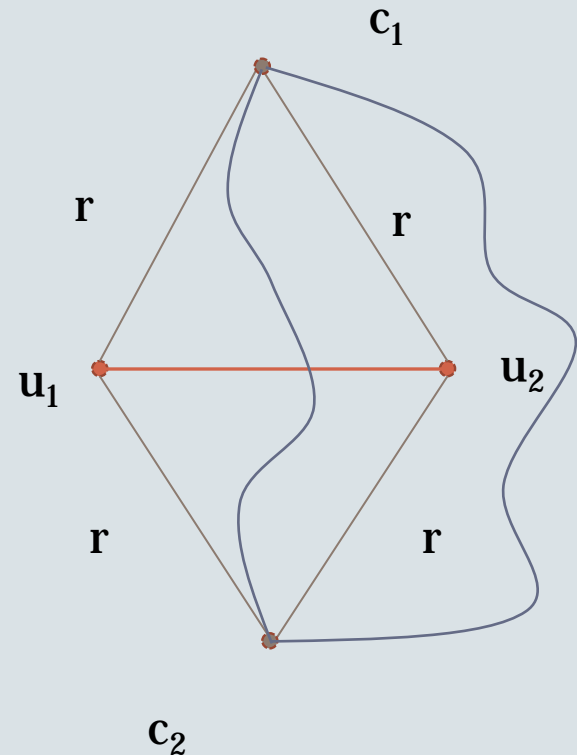
1. we don't define this  $N_G(u_1, u_2, \beta)$  set for  $\beta=0$  nor for  $\beta=\infty$

We will explain in a moment why.

2. for  $0 < \beta < 1$  let points  $c_1$  and  $c_2$  be such points (if they exist) that:  
 $d_G(c_1, u_1) = d_G(c_2, u_1) = d_G(c_1, u_2) = d_G(c_2, u_2) =$   
 $= d_G(u_1, u_2) / 2\beta$  and all paths connecting those points with length shorter than  $d_G(u_1, u_2) / \beta$  intersect shortest paths between  $u_1$  and  $u_2$ :

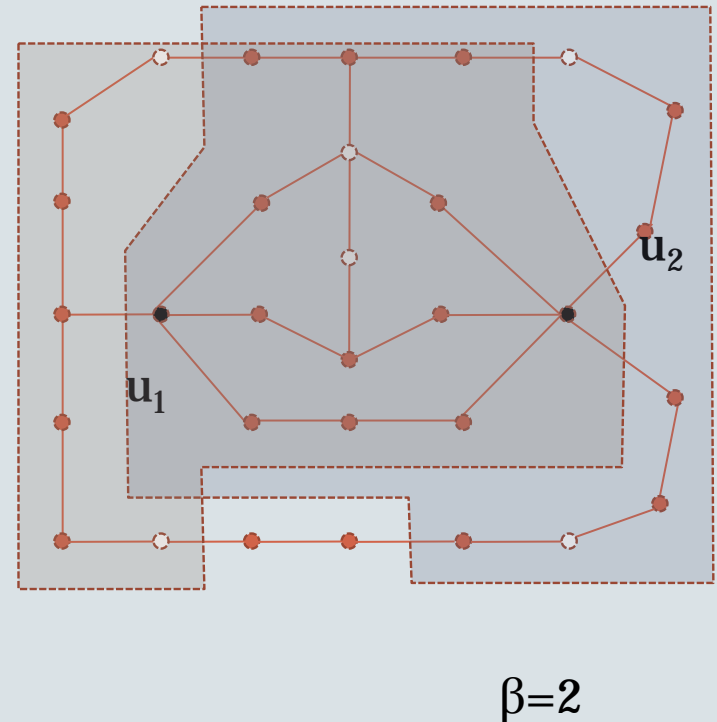
then  $N_G(u_1, u_2, \beta)$  is an intersection of discs  $D_G(c_1, d_G(u_1, u_2) / 2\beta)$  and  $D_G(c_2, d_G(u_1, u_2) / 2\beta)$ ;

if points  $c_1$  and  $c_2$  don't exist then  $N_G(u_1, u_2, \beta)$  is an empty set;



$$d(u_1, u_2) / 2\beta = r$$

3. for  $1 \leq \beta < \infty$  we define set  $C_1$  (respectively  $C_2$ ) of disc centers  $c_1$  such that the distance between  $u_1$  and  $c_1$  ( $u_2$  and  $c_2$  respectively) is  $\beta d_G(u_1, u_2)/2$  and the distance between  $u_1$  and  $c_2$  ( $u_2$  and  $c_1$  respectively) is  $|1 - \beta/2| d_G(u_1, u_2)$ ; set  $N_G(u_1, u_2, \beta)$  is an intersection of any two discs in  $G$  each with radius  $\beta d_G(u_1, u_2)/2$ , centered in such points  $c_1 \in C_1$  and  $c_2 \in C_2$  (respectively) tht the distance between  $c_1$  and  $c_2$  is  $(\beta - 1) |v_1 v_2|$ ;



We call graph  $G_\beta(G)$  a **lune-based  $\beta$ -skeleton** for graph  $G$ .

## Lemma

For each graph  $G=(V,U,E)$  and for  $1 \leq \beta \leq \beta' \leq 2$  following inclusions are true:

$$\text{MST}(G) \subseteq \text{RNG}(G) \subseteq G_{\beta'}(G) \subseteq G_{\beta}(G) \subseteq \text{GG}(G) \subseteq \text{DT}(G)$$

where  $\text{RNG}(G) = G_2(G)$  and  $\text{GG}(G) = G_1(G)$ .

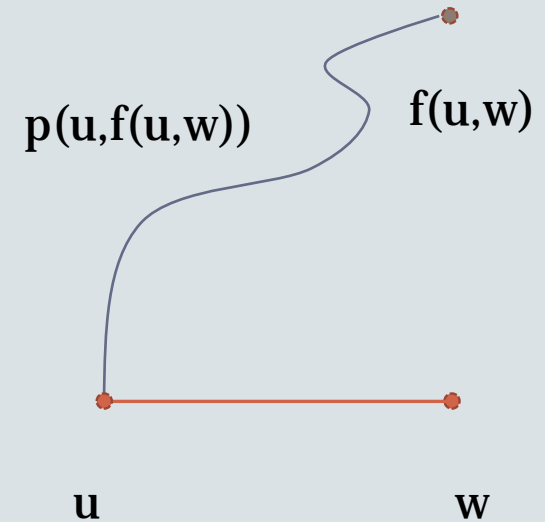
We cannot modify definition of the lune-based beta skeleton for  $\beta=0$  and for  $\beta=\infty$  because if the weights of the edges of the graph  $G$  are finite then we can only define  $\beta$ -skeletons for some values of  $\beta$ .

## Lemma

For graph  $G=(V,U,E)$  we can at most define  $G_\beta(G)$  for  $\beta_{\min} \leq \beta \leq \beta_{\max}$  where:

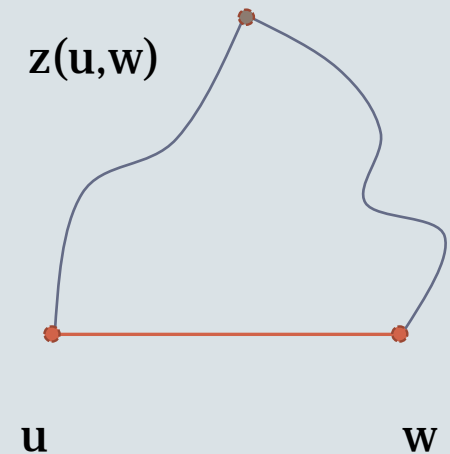
1. For a given pair of vertices  $u, w$  in graph  $G$  let  $f(u,w)$  be the point farthest from  $u$  in graph  $G \setminus \{\text{all edges on shortest paths between } u \text{ and } w\}$  and let  $p(u, f(u,w))$  be the shortest path between  $u$  and  $f(u,w)$  in this graph.

Then,  $\beta_{\max} = 2[(\max_{u,w} |p(u, f(u,w))| / |uw|)] + 2$ .



2. Now, let  $z(u, w)$  be the farthest point from  $u_1$  and  $u_2$  such that  $d(u, z(u, w)) = d(w, z(u, w))$  in graph  $G \setminus \{\text{all edges on shortest paths between } u \text{ and } w\}$ .

Then,  $\beta_{\min} = \min_{u,w} [|uw| / 2(d(u, z(u, w)))]$ .



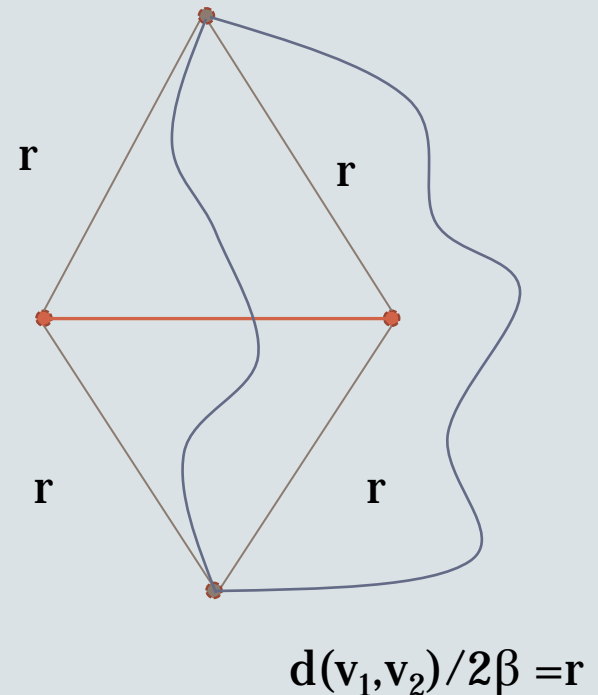


# Generalized $\beta$ -skeleton

25

For a given **set of objects**  $S$  in space  $L$  with metric  $d$  we define a graph  $G_\beta(S)$ , where an edge exists between two points  $v_1$  from  $s_1$  and  $v_2$  from  $s_2$  iff at least one set  $N_d(v_1, v_2, \beta)$  is not intersected by any object from  $S \setminus \{s_1, s_2\}$  where:

1. for  $0 < \beta < 1$  let points  $c_1$  and  $c_2$  be such points (if they exist) that:  $d(c_1, v_1) = d(c_2, v_1) = d(c_1, v_2) = d(c_2, v_2)$  and all paths connecting those points with length shorter than  $d(v_1, v_2)/2\beta$  intersect shortest paths between  $v_1$  and  $v_2$ : then  $N_d(u_1, u_2, \beta)$  is an intersection of discs  $D(c_1, d(v_1, v_2)/2\beta)$  and  $D(c_2, d(v_1, v_2)/2\beta)$ ;  
if points  $c_1$  and  $c_2$  don't exist then  $N_d(v_1, v_2, \beta)$  is an empty set;



2. for  $1 \leq \beta < \infty$  we define set  $C_1$  (respectively  $C_2$ ) of disc centers  $c_1$  such that the distance between  $v_1$  and  $c_1$  ( $v_2$  and  $c_2$  respectively) is  $\beta d(v_1, v_2)/2$  and the distance between  $v_1$  and  $c_2$  ( $v_2$  and  $c_1$  respectively) is  $|1 - \beta/2| d(v_1, v_2)$ ; set  $N_d(v_1, v_2, \beta)$  is an intersection of any two discs in  $L$ , each with radius  $\beta d(v_1, v_2)/2$ , centered in such points  $c_1 \in C_1$  and  $c_2 \in C_2$  that the distance between  $c_1$  and  $c_2$  is  $(\beta - 1) |v_1 v_2|$ ;

3. for  $\beta=0$  let  $C_0$  be a set of all discs  $c$  such that there exists a sequence of discs  $\{c(\beta) \mid c(\beta) \text{ is a disc that we used to define set } N_d(v_1, v_2, \beta) \}$  convergent to  $c$  with  $\beta \rightarrow 0$ ;  
 $N_d(v_1, v_2, 0)$  is an intersection of any two different discs from  $C_0$ ;
  
4. for  $\beta=\infty$  we define  $C_\infty$  as a set of all discs  $c$  such that there exists a sequence of discs  $\{c(\beta) \mid c(\beta) \text{ is a disc that we used to define set } N_d(v_1, v_2, \beta) \}$  convergent to  $c$  with  $\beta \rightarrow \infty$ ;  
then  $N_d(v_1, v_2, \infty)$  is intersection of any two different discs from  $C_\infty$ .

# Conclusions and open problems

28

We showed a way to generalize  $\beta$ -skeletons basing on distance criterion. We focused only on few special cases but we think that they describe well the idea if this general definition.

In a similar way we could define  $\beta$ -skeletons for sets of polygons and it is also possible to generalize this definition to higher dimensions.

There is a couple of new problems regarding this definition. It would be interesting to check how those changes can influence the time of algorithms computing  $\beta$ -skeletons.

We can also analyse what interesting properties do  $\beta$ -skeletons for different objects have.

**Thank You for Your Attention**